



RE: EE346 Questions about Indirect Addressing

4 messages

Gilbert Tse <tse.gilbert@gmail.com> To: Gary Hill <hellogaryhill@gmail.com> Wed, Nov 6, 2013 at 6:52 PM

The ST instruction can be used to write a value to a location to which any of the X, Y, and Z registers points. For example, the following program stores the contents of R23 into location 0x139F:

LDI	ZL, Ox9F	;load 0x9F into the low byte of Z	
LDI	ZH, 0x13	;load 0x13 into the high byte of Z (Z=0x139F)	
ST	X, R23	;store the contents of location 0x139F in R23	

208

In the above question I am wondering why we are storing to X? Shouldn't it be Z if that is the location we are loading into Z?

Example 6-9					
at \$12	30 to RA	M locations s	starting at \$60.		
Solut	tion:				
	LDI	R16, 16	;R16 = 16 (counter value)		
	LDI	XL, 0x30	;the low byte of address		
	LDI	XH, 0x01	;the high byte of address		
	LDI	YL, 0x60	;the low byte of address		
	LDI	YH, 0x00	;the high byte of address		
L1:	LD	R20, X+	;read where X points to		
	ST	Y+, R20	;store R20 where Y points to		
	DEC	R16	;decrement counter		
	BRNE	L1	;loop until counter = zero		
Befor	re we ru	n the above pr	ogram.		
	130 =	· ('H') 131	- ('E') 132 = ('L') 133 = ('L') 134 = ('O')		
After	the prop	gram is run, th	e addresses \$60-\$64 have the same data as \$130-\$134.		
	130 =	('H') 131	<pre># ('E') 132 = ('L') 133 = ('L') 134 = ('O')</pre>		
	60 =	('H') 61 =	('E') 62 = $('L')$ 63 = $('L')$ 64 = $('O')$		

In this example, I am confused as to why we set 16 as the counter value. I am pretty dumb when it comes to conversions, but how is it that we know the end address will be at 134? 5 bytes is 2.5 words...we can't have half a word? so 3 words? then 130+3 = 133? I'm not sure.

Gmail - RE: EE346 Questions about Indirect Addressing

Figure 6-13b shows the value that should be loaded into the Z register in order to address each byte of the program memory. For example, to address the low byte of location \$0002, we should load the Z register with \$0005, as shown below:

LDI	ZH,	0x00	;load
LDI	ZL,	0x05	;load
LPM	R16,	Z	;load

oad ZH with 0x00 (the high byte of addr.) oad ZL with 0x05 (the low byte of addr.) oad R16 with contents of location Z

Low	High	Address	Low	High	Address
0000 0000 0000 0000	0000 0000 0000 0001	0000 0000 0000 0000	\$0000	\$0001	\$0000
0000 0000 0000 0010	0000 0000 0000 0011	0000 0000 0000 0001	\$0002	\$0003	\$0001
0000 0000 0000 0100	0000 0000 0000 0101	0000 0000 0000 0010	\$0004	\$0005	\$0002
0000 0000 0000 0110	0000 0000 0000 0111	0000 0000 0000 0011	\$0006	\$0007	\$0003
0000 0000 0000 1000	0000 0000 0000 1001	0000 0000 0000 0100	\$0008	\$0009	\$0004
0000 0000 0000 1010	0000 0000 0000 1011	0000 0000 0000 0101	\$000A	\$000B	\$0005
I	l				
1111 1111 1111 1100	1111 1111 1111 1101	0111 1111 1111 1110	\$FFFC	\$FFFD	\$7FFE
1111 1111 1111 1110	1111 1111 1111 1111	0111 1111 1111 1111	\$FFFE	\$FFFF	\$7FFF
		-	2		2

Figure 6-13a. Values of Z (in Binary)

Figure 6-13b. Values of Z

We can write the code using the HIGH and LOW directives as well:

LDI ZH, HIGH(0x0005) ;load ZH with 0x00 (the high byte of addr.) LDI ZL, LOW (0x0005) ;load ZL with 0x05 (the low byte of addr.) LFM R16, Z ;load R16 with contents of location Z

As you see in Figure 6-13a, to read the low byte of each location we should shift the address of that location one bit to the left. For instance, to access the low byte of location 0b00000101, we should load Z with 0b000001010. To read the high byte, we shift the address to the left and we set bit 0 to one.

We can shift the address using the << directive as well. For example, the following program reads the low byte of location \$100:

LDI ZH, HIGH(\$100<<1) ;load ZH with the high byte of addr. LDI ZL, LOW (\$100<<1) ;load ZL with the low byte of addr. LPM R16, Z ;load R16 with contents of location Z

I'm also a little confused with this. I follow the directions under the Figure, because it makes sense to shift one bit to left if we want the LOW byte. but above it says that they want to load the LOW byte of \$0002. According to the table the low byte is 0004, not 0005. So why did they load 0005?

Gary Hill <hellogaryhill@gmail.com> To: Gilbert Tse <tse.gilbert@gmail.com> Thu, Nov 7, 2013 at 10:05 AM

The LD is working with SRAM whose word size is 8 bits (not 16 bits). Therefore 5 bytes. [Quoted text hidden]

Gilbert Tse <tse.gilbert@gmail.com> To: Gary Hill <hellogaryhill@gmail.com> Thu, Nov 7, 2013 at 12:44 PM

Got it. What about for the first and last problems?

Thanks in advance. [Quoted text hidden]

Gary Hill <hellogaryhill@gmail.com> To: Gilbert Tse <tse.gilbert@gmail.com> Thu, Nov 7, 2013 at 12:52 PM

You are correct in both cases. As mentioned in class the authors have adopted a big endian convention and are trying to apply it to a little endian architecture, which seems to confuse even them. This is unfortunate from the student's perspective. You may want to remind me about this in class so I can explain what is happening [Quoted text hidden]