## RE: EE346 Questions about Indirect Addressing

4 messages
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Wed, Nov 6, 2013 at 6:52 PM
To: Gary Hill [hellogaryhill@gmail.com](mailto:hellogaryhill@gmail.com)
The ST instruction can be used to write a value to a location to which any of the $\mathrm{X}, \mathrm{Y}$, and Z registers points. For example, the following program stores the contents of R23 into location 0x139F:

| LDI | ZL, $0 \times 9 \mathrm{E}$ | ; load $0 \times 9 \mathrm{E}$ into the low byte of Z |
| :--- | :--- | :--- |
| LDI | $\mathrm{ZH}, 0 \times 13$ | ;load $0 \times 13$ into the high byte of $\mathrm{Z}(2=0 \times 139 \mathrm{~F})$ |
| ST | $\mathrm{X}, \mathrm{R} 23$ | istore the contents of location $0 \times 139 \mathrm{~F}$ in R23 |

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In the above question I am wondering why we are storing to $X$ ? Shouldn't it be $Z$ if that is the location we are loading into Z ?

## Example 6-9

Write a program to copy a block of 5 bytes of data from data memory locations starting at $\$ 130$ to RAM locations starting at $\$ 60$.

Solution:

| L1: | LDI | R16, 16 | ;R16 - 16 (counter value) |
| :---: | :---: | :---: | :---: |
|  | LDI | XL, 0x30 | ; the low byte of address |
|  | LDI | XH, 0x01 | ; the high byte of address |
|  | LDI | YL, 0x60 | ; the low byte of address |
|  | LDI | YH, $0 \times 00$ | ; the high byte of address |
|  | LD | R20, X+ | ; read where X points to |
|  | ST | $\mathrm{Y}+$, R20 | ;store R20 where $Y$ points to |
|  | DEC | R16 | ; decrement counter |
|  | BRNE | L1 | ; loop until counter = zero |

Before we run the above program.

```
130=('H') 131 = ('E') 132 = ('L') 133 = ('L') 134 = ('O')
```

After the program is run, the addresses $\$ 60-\$ 64$ have the same data as $\$ 130-\$ 134$.

```
130 = ('H') 131 m ('E') 132 = ('L') 133 = ('L') 134 = ('O')
60=('H') 61 = ('E') 62 = ('L') 63 = ('L') 64 = ('O')
```

In this example, I am confused as to why we set 16 as the counter value. I am pretty dumb when it comes to conversions, but how is it that we know the end address will be at 134 ? 5 bytes is 2.5 words...we can't have half a word? so 3 words? then $130+3=133$ ? I'm not sure.

Figure 6-13b shows the value that should be loaded into the Z register in order to address each byte of the program memory. For example, to address the low byte of location $\$ 0002$, we should load the Z register with $\$ 0005$, as shown below:

$$
\begin{array}{ll}
\text { LDI } 2 H, 0 \times 00 & \text { iload } 2 H \text { with } 0 \times 00 \text { (the high byte of addr.) } \\
\text { LDI } 2 L, ~ 0 x 05 & \text { iload } 2 I \text { with } 0 \times 05 \text { (the low byte of addr.) } \\
\text { LPM RIG, } Z & \text { iload R16 with contents of location } 2
\end{array}
$$



Figure 6-13a. Values of $Z$ (in Binary)

| Low | High | Address <br> $\$ 0000$ |
| :---: | :---: | :---: |
| \$0000 | \$0001 |  |
| \$0002 | \$0003 | \$0001 |
| \$0004 | \$0005 | \$0002 |
| \$0006 | \$0007 | \$0003 |
| \$0008 | \$0009 | \$0004 |
| \$000A | \$000B | \$0005 |
| \$FFFC | \$FFFD | \$7FFE |
| \$FFFE | \$FFFF | S7FFF |

Figure 6-13b. Values of $Z$

We can write the code using the HIGH and LOW directives as well:

```
LDI ZH, HIGH (OxOOO5) ; load ZH with 0x00 (the high byte of addr.)
LDI ZL, IOW (0x0005) ; load zL with 0x05 (the low byte of addz.)
LPM R16, }Z\mathrm{ ;load R16 with contents of location }
```

As you see in Figure 6-13a, to read the low byte of each location we should shift the address of that location one bit to the left. For instance, to access the low byte of location 0b00000101, we should load Z with 0 b 000001010 . To read the high byte, we shift the address to the left and we set bit 0 to one.

We can shift the address using the $\ll$ directive as well. For example, the following program reads the low byte of location $\$ 100$ :

```
LDI ZH, BIGH($100<<1) ; load ZH with the high byte of adcir.
LDI 2L, LOW ($100<<1) ; load ZL with the low byte of addr.
LPM RlG, 2 ; load Rl6 with contents of location Z
```

I'm also a little confused with this. I follow the directions under the Figure, because it makes sense to shift one bit to left if we want the LOW byte. but above it says that they want to load the LOW byte of $\$ 0002$. According to the table the low byte is 0004, not 0005 . So why did they load 0005 ?

Gary Hill [hellogaryhill@gmail.com](mailto:hellogaryhill@gmail.com)
Thu, Nov 7, 2013 at 10:05 AM
To: Gilbert Tse [tse.gilbert@gmail.com](mailto:tse.gilbert@gmail.com)
The LD is working with SRAM whose word size is 8 bits (not 16 bits). Therefore 5 bytes.
[Quoted text hidden]

Gilbert Tse [tse.gilbert@gmail.com](mailto:tse.gilbert@gmail.com)
Thu, Nov 7, 2013 at 12:44 PM
To: Gary Hill [hellogaryhill@gmail.com](mailto:hellogaryhill@gmail.com)

Got it. What about for the first and last problems?
Thanks in advance.
[Quoted text hidden]

Gary Hill [hellogaryhill@gmail.com](mailto:hellogaryhill@gmail.com)
Thu, Nov 7, 2013 at 12:52 PM
To: Gilbert Tse [tse.gilbert@gmail.com](mailto:tse.gilbert@gmail.com)
You are correct in both cases. As mentioned in class the authors have adopted a big endian convention and are trying to apply it to a little endian architecture, which seems to confuse even them. This is unfortunate from the student's perspective. You may want to remind me about this in class so I can explain what is happening
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